

C^* -algebras of graph products

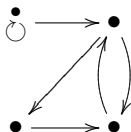
Ann Johnston (Harvey Mudd College)
Andrew Reynolds (UCLA)

July 22, 2009

Definition of a (directed) Graph

A graph $E = (E^0, E^1, r, s)$ consists of two countable sets

- ▶ E^0 : vertices
- ▶ E^1 : edges
- ▶ and functions $r, s : E^1 \rightarrow E^0$ called *range and source*



Definition of a C^* -algebra

Let F be a field. A vector space A over F with a binary operation $A \times A \rightarrow A$ (called multiplication) is an *associative algebra over F* if $\forall x, y \in A$ and $\forall a, b \in F$

- ▶ $(x + y) \cdot z = x \cdot z + y \cdot z$
- ▶ $x \cdot (y + z) = x \cdot y + x \cdot z$
- ▶ $(ax) \cdot (by) = (ab)(x \cdot y)$
- ▶ $(x \cdot y) \cdot z = x \cdot (y \cdot z)$.

Definition of a C^* -algebra

Let A be an associative algebra over \mathbb{C} .

A *norm* on A is a map $\|\cdot\| : A \rightarrow \mathbb{R}$ satisfying

- ▶ $\|v\| \geq 0$ for all $v \in A$ and $v = 0$ if and only if $v = 0$ in A
- ▶ $\|\lambda v\| = |\lambda| \|v\|$ for all $v \in A$ and all $\lambda \in \mathbb{C}$
- ▶ $\|v + w\| \leq \|v\| + \|w\|$ for all $v, w \in A$.

An *involution* on a complex algebra A is a real-linear map $T \mapsto T^*$ such that for all $S, T \in A$ and $\lambda \in \mathbb{C}$, we have

$$T^{**} = T, \quad (ST)^* = T^*S^*, \quad (\lambda T)^* = \bar{\lambda}T^*.$$

Definition of a C^* -algebra

A C^* -algebra A is an associative algebra over \mathbb{C} with a norm $a \mapsto \|a\|$ and an involution $a \mapsto a^*$, such that A is complete with respect to the norm, and such that $\|ab\| \leq \|a\|\|b\|$ and $\|a^*a\| = \|a\|^2$ for every $a, b \in A$.

Examples of C^* -algebras

In general a C^* -algebra can be infinite dimensional. It is a theorem (Gelfand-Naimark) that every C^* -algebra is isomorphic to a sub C^* -algebra of $B(H)$, the algebra of all bounded linear operators on some Hilbert space H .

Our examples will be limited to finite dimensional case, and so our C^* -algebras can be thought of as $M_n(\mathbb{C})$, with involution defined as taking the conjugate transpose, and the norm being the operator norm on matrices:

$$\|A\|_{op} = \sup_{\|x\| \leq 1} \|Ax\| = \sup_{\|x\|=1} \|Ax\|.$$

Forming the C^* -algebra of a graph

Let E be a row-finite directed graph. A *Cuntz-Krieger E -family* $\{S, P\}$ on a Hilbert space H consists of a set $\{P_v : v \in E^0\}$ of mutually orthogonal projections on H , and a set $\{S_e : e \in E^1\}$ of partial isometries on H such that

- ▶ $S_e^* S_e = P_{s(e)} \forall e \in E^1$; and
- ▶ $P_v = \sum_{\{e \in E^1 : r(e)=v\}} S_e S_e^*$ whenever v is not a source.

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It is a theorem that for E there is a C^* -algebra, $C^*(E)$ generated by a Cuntz-Krieger E -family which has a universal property. It is unique up to isomorphism and is called the *graph algebra* of E .

Note: Since we are dealing only with finite-dimensional cases, we will always have $H \cong \mathbb{C}^n$ for some n .

Example

$$E : a \xrightarrow{e} b$$

► Relations:

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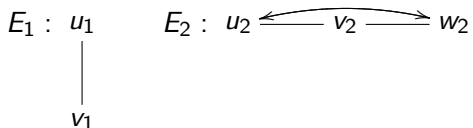
- $S_e^* S_e = P_a$
- $P_b = S_e S_e^*$

► Generators:

$$P_a = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, \quad P_b = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}, \quad S_e = \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}$$

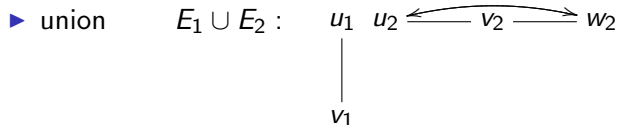
- $C^*(E) = C^*(P_a, P_b, S_e) = M_2(\mathbb{C})$.

Products of graphs

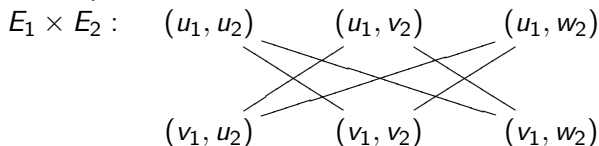


There are many ways of forming larger graphs given two smaller graphs. Examples are:

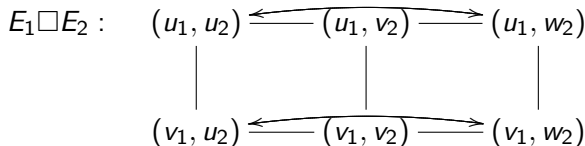
Products of graphs



- tensor product



- box (cartesian) product



Products of graphs

The *box product* of E_1 with E_2 is the graph $E_1 \square E_2$ with vertex set $E_1^0 \times E_2^0$, where for all $(u_1, u_2), (v_1, v_2) \in E_1^0 \times E_2^0$ we define $(u_1, u_2) \sim (v_1, v_2)$ if either one of the following holds :

- ▶ $u_1 = v_1$ in G_1 and $u_2 \sim v_2$ in G_2 ,
- ▶ $u_2 = v_2$ in G_2 and $u_1 \sim v_1$ in G_1 .

$$E : \quad a \xrightarrow{e} b, \quad E \square E : \quad \begin{array}{ccc} (a, a) & \longrightarrow & (a, b) \\ \downarrow & & \downarrow \\ (b, a) & \longrightarrow & (b, b) \end{array}$$

C*-algebra of the box product

$$E : a \xrightarrow{e} b,$$

$E \square E :$

$$\begin{array}{ccc} w & \xrightarrow{e} & x \\ g \downarrow & & \downarrow h \\ y & \xrightarrow{f} & z \end{array}$$

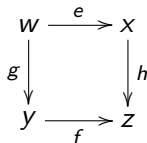
► Relations:

- $S_e^* S_e = P_w = S_g^* S_g$
- $S_h^* S_h = P_x = S_e S_e^*$
- $S_f^* S_f = P_y = S_g S_g^*$
- $P_z = S_f S_f^* + S_h S_h^*$

C^* -algebra of the box product

$$E : a \xrightarrow{e} b,$$

$$E \square E :$$



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- $S_e^* S_e = P_w = S_g^* S_g$
 - $S_h^* S_h = P_x = S_e S_e^*$
 - $S_f^* S_f = P_y = S_g S_g^*$
 - $P_z = S_f S_f^* + S_h S_h^*$
- We see that this requires a five-dimensional hilbert space, and with calculations we find the graph algebra to be isomorphic to $M_5(\mathbb{C})$.
- We have seen that $C^*(E) \cong M_2(\mathbb{C})$, and so we see that $C^*(E) \otimes C^*(E) = M_4(\mathbb{C}) \not\cong C^*(E \square E)$.

2-graphs

A *category* C consists of two classes C^0 of objects, and C^* of morphisms, and two functions $r, s : C^* \rightarrow C^0$, domain and codomain, as well as a partially defined product (composition) $(f, g) \mapsto f \circ g$ from $\{(f, g) \in C^* \times C^* : s(f) = r(g)\}$ to C^* , composition, and distinguished elements (identity morphisms) $\{i_v \in C^* : v \in C^0\}$ satisfying

- ▶ $r(fg) = r(f)$ and $s(fg) = s(g)$
- ▶ $(fg)h = f(gh)$ when $s(f) = r(g)$ and $s(g) = r(h)$
- ▶ $r(i_v) = v = s(i_v)$ and $i_v f = f, g i_v = g$ when $r(f) = v$ and $s(g) = v$.

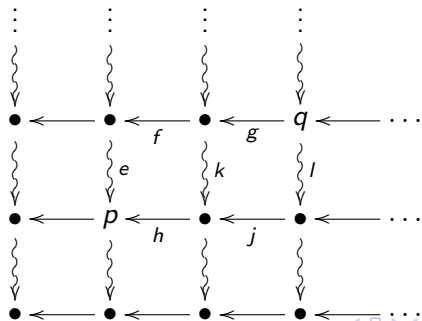
A *functor* $F : C \rightarrow D$ is a pair of maps $F^0 : C^0 \rightarrow D^0$ and $F^* : C^* \rightarrow D^*$ which respect the domain and codomain maps and composition, and which satisfy $F^*(i_v) = i_{F^0(v)}$.

2-graphs

A *rank 2 graph* (or *2-graph*) (Λ, d) is a countable category Λ , together with a functor $d : \Lambda \rightarrow \mathbb{N}^2$, called the degree map, with the following unique factorization property:

for every morphism λ and every decomposition $d(\lambda) = m + n$ with $m, n \in \mathbb{N}^2$, there exist unique morphisms μ and ν such that $d(\mu) = m$, $d(\nu) = n$ and $\lambda = \mu\nu$.

For every 2-graph Λ , there is an associated 1-skeleton E_Λ , which is a colored directed graph with $E^0 = \Lambda^0$ and $E^1 = \Lambda^{(0,1)} \cup \Lambda^{(1,0)}$.



2-graphs

Given categories C and D , the *product category* $C \times D$ is the category with

- ▶ $(C \times D)^0 = C^0 \times D^0$
- ▶ $(C \times D)^* = C^* \times D^*$, with composition component-wise from the contributing categories and identities $(1_A, 1_B)$ where $A \in C^0, B \in D^0$.

Forming a 2-graph from two graphs

Viewing two graphs, E and F , as categories with objects as vertices and morphisms as paths, we can form their product category, which we then view as a 2-graph $E \times F$, with the degree map being ordered pairs of path lengths.

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Viewing two graphs, E and F , as categories with objects as vertices and morphisms as paths, we can form their product category, which we then view as a 2-graph $E \times F$, with the degree map being ordered pairs of path lengths.

Example:

$$E : \quad a \xrightarrow{e} b \quad \text{1-skeleton for } E \times E : \quad \begin{array}{ccc} w & \xrightarrow{e} & x \\ \left. \begin{array}{c} \text{\scriptsize } g \\ \text{\scriptsize } \downarrow \\ \text{\scriptsize } y \end{array} \right\} & & \left. \begin{array}{c} \text{\scriptsize } h \\ \text{\scriptsize } \downarrow \\ \text{\scriptsize } z \end{array} \right\} \\ y & \xrightarrow{f} & z \end{array}$$

Notice that ignoring color, $E \square E$ is isomorphic to the 1-skeleton of $E \times E$.

C^* -algebras of 2-graphs

$$E : \quad a \xrightarrow{e} b \quad \text{1-skeleton for } E \times E : \quad \begin{array}{ccc} w & \xrightarrow{e} & x \\ \left. \begin{array}{c} \text{\scriptsize } g \\ \text{\scriptsize } \downarrow \\ \text{\scriptsize } y \end{array} \right\} & & \left. \begin{array}{c} \text{\scriptsize } h \\ \text{\scriptsize } \downarrow \\ \text{\scriptsize } z \end{array} \right\} \\ & \xrightarrow{f} & \end{array}$$

There are slightly modified Cuntz-Krieger relations for 2-graphs, which account for the degrees of the morphisms of the graph. Using these relations, we still generate a universal graph algebra unique up to isomorphism. Under these relations $C^*(E \times E) \cong M_4(\mathbb{C}) \cong C^*(E) \otimes C^*(E)$.

In general, Alex Kumjian and David Pask have shown that the C^* -algebra of the 2-graph constructed from two graphs, E, F is isomorphic to the tensor product of their corresponding C^* -algebras. That is,

$$C^*(E \times F) \cong C^*(E) \otimes C^*(F).$$

Work in progress

We are in the process of investigating several other products. We have already investigated

- ▶ direct sum of algebras
- ▶ crossed product of an algebra with a group

We are still considering:

- ▶ tensor product of graphs
- ▶ an “overlay” product, that would lead to the strong product
- ▶ etc.

References

- [1] Kenneth R. Davidson, “C*-Algebras by Example”. Providence: American Mathematical Society, 1996.
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- [3] Iain Raeburn, “Graph Algebras”. Providence: American Mathematical Society, 2005.
- [4] M. Rordam, F. Larsen, & N. J. Laustsen, “An Introduction to K-Theory for C*-Algebras”. Cambridge: Cambridge University Press, 2000.